

Preimages of HAVAL and MD5

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Definition

“Given y , find x such that $h(x) = y$ ”

Given what?

1. **random range** element
2. **random image**
3. **any fixed** value (“weak” images)

Attacks here work for 1, 2, 3

Previous preimage attacks

Google (*collision attack*) $\approx 500\,000$

Google (*preimage attack*) $\approx 15\,000$

- ▶ MD2 (2004)
- ▶ Parallel FFT-Hashing (2007)
- ▶ Syndrome-based hash (2007)
- ▶ HAS-V (2007)
- ▶ Tiger (2007)
- ▶ MD4 (2008)
- ▶ GOST (2008)
- ▶ Snejfru (2008)
- ▶ SHA-0/1 (2008)

Content of this talk

MD5: how to invert...

- ▶ 22 steps?
- ▶ 31 steps?
- ▶ 47 steps?

What about **HAVAL**?

From the **compression to hash** function

MD5

- ▶ 1991: publication (Rivest)
- ▶ 1993: **collision** attack (compression function)
- ▶ 2005: **collision** attack (hash function)
- ▶ 2005+: faster, chosen-prefix, meaningful **collisions**

No preimage attack.



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MD5 compression function

Input

- ▶ chain value $H_0 H_1 H_2 H_3$
- ▶ message $M_0 M_1 \dots M_{15}$

Algorithm

- ▶ copy $H_0 H_1 H_2 H_3$ into $A_0 B_0 C_0 D_0$
- ▶ for $i = 1 \dots 64$

$$\begin{aligned} A_i &= D_{i-1} \\ B_i &= f_i(A_{i-1}, B_{i-1}, C_{i-1}, D_{i-1}, M_{\sigma(i)}) \\ C_i &= B_{i-1} \\ D_i &= C_{i-1} \end{aligned}$$

- ▶ return $(A_0 + A_{64}) \parallel (B_0 + B_{64}) \parallel (C_0 + C_{64}) \parallel (D_0 + D_{64})$

Unrolling (index, function, message word)

1	$f_1(A_0, B_0, C_0, D_0, 0)$	17	$f_{17}(A_{16}, B_{16}, C_{16}, D_{16}, 1)$
2	$f_2(A_1, B_1, C_1, D_1, 1)$	18	$f_{18}(A_{17}, B_{17}, C_{17}, D_{17}, 6)$
3	$f_3(A_2, B_2, C_2, D_2, 2)$	19	$f_{19}(A_{18}, B_{18}, C_{18}, D_{18}, 11)$
4	$f_4(A_3, B_3, C_3, D_3, 3)$	20	$f_{20}(A_{19}, B_{19}, C_{19}, D_{19}, 0)$
5	$f_5(A_4, B_4, C_4, D_4, 4)$	21	$f_{21}(A_{20}, B_{20}, C_{20}, D_{20}, 5)$
6	$f_6(A_5, B_5, C_5, D_5, 5)$	22	$f_{22}(A_{21}, B_{21}, C_{21}, D_{21}, 10)$
7	$f_7(A_6, B_6, C_6, D_6, 6)$	23	$f_{23}(A_{22}, B_{22}, C_{22}, D_{22}, 15)$
8	$f_8(A_7, B_7, C_7, D_7, 7)$	24	$f_{24}(A_{23}, B_{23}, C_{23}, D_{23}, 4)$
9	$f_9(A_8, B_8, C_8, D_8, 8)$	25	$f_{25}(A_{24}, B_{24}, C_{24}, D_{24}, 9)$
10	$f_{10}(A_9, B_9, C_9, D_9, 9)$	26	$f_{26}(A_{25}, B_{25}, C_{25}, D_{25}, 14)$
11	$f_{11}(A_{10}, B_{10}, C_{10}, D_{10}, 10)$	27	$f_{27}(A_{26}, B_{26}, C_{26}, D_{26}, 3)$
12	$f_{12}(A_{11}, B_{11}, C_{11}, D_{11}, 11)$	28	$f_{28}(A_{27}, B_{27}, C_{27}, D_{27}, 8)$
13	$f_{13}(A_{12}, B_{12}, C_{12}, D_{12}, 12)$	29	$f_{29}(A_{28}, B_{28}, C_{28}, D_{28}, 13)$
14	$f_{14}(A_{13}, B_{13}, C_{13}, D_{13}, 13)$	30	$f_{30}(A_{29}, B_{29}, C_{29}, D_{29}, 2)$
15	$f_{15}(A_{14}, B_{14}, C_{14}, D_{14}, 14)$	31	$f_{31}(A_{30}, B_{30}, C_{30}, D_{30}, 7)$
16	$f_{16}(A_{15}, B_{15}, C_{15}, D_{15}, 15)$	32	$f_{32}(A_{31}, B_{31}, C_{31}, D_{31}, 12)$

First 2 rounds

- | | | | |
|----|---|----|---|
| 1 | $f(A_0, B_0, C_0, D_0, 0)$ | 17 | $g(A_{16}, B_{16}, C_{16}, D_{16}, 1)$ |
| 2 | $f(A_1, B_1, C_1, D_1, 1)$ | 18 | $g(A_{17}, B_{17}, C_{17}, D_{17}, 6)$ |
| 3 | $f(A_2, B_2, C_2, D_2, 2)$ | 19 | $g(A_{18}, B_{18}, C_{18}, D_{18}, 11)$ |
| 4 | $f(A_3, B_3, C_3, D_3, 3)$ | 20 | $g(A_{19}, B_{19}, C_{19}, D_{19}, 0)$ |
| 5 | $f(A_4, B_4, C_4, D_4, 4)$ | 21 | $g(A_{20}, B_{20}, C_{20}, D_{20}, 5)$ |
| 6 | $f(A_5, B_5, C_5, D_5, 5)$ | 22 | $g(A_{21}, B_{21}, C_{21}, D_{21}, 10)$ |
| 7 | $f(A_6, B_6, C_6, D_6, 6)$ | 23 | $g(A_{22}, B_{22}, C_{22}, D_{22}, 15)$ |
| 8 | $f(A_7, B_7, C_7, D_7, 7)$ | 24 | $g(A_{23}, B_{23}, C_{23}, D_{23}, 4)$ |
| 9 | $f(A_8, B_8, C_8, D_8, 8)$ | 25 | $g(A_{24}, B_{24}, C_{24}, D_{24}, 9)$ |
| 10 | $f(A_9, B_9, C_9, D_9, 9)$ | 26 | $g(A_{25}, B_{25}, C_{25}, D_{25}, 14)$ |
| 11 | $f(A_{10}, B_{10}, C_{10}, D_{10}, 10)$ | 27 | $g(A_{26}, B_{26}, C_{26}, D_{26}, 3)$ |
| 12 | $f(A_{11}, B_{11}, C_{11}, D_{11}, 11)$ | 28 | $g(A_{27}, B_{27}, C_{27}, D_{27}, 8)$ |
| 13 | $f(A_{12}, B_{12}, C_{12}, D_{12}, 12)$ | 29 | $g(A_{28}, B_{28}, C_{28}, D_{28}, 13)$ |
| 14 | $f(A_{13}, B_{13}, C_{13}, D_{13}, 13)$ | 30 | $g(A_{29}, B_{29}, C_{29}, D_{29}, 2)$ |
| 15 | $f(A_{14}, B_{14}, C_{14}, D_{14}, 14)$ | 31 | $g(A_{30}, B_{30}, C_{30}, D_{30}, 7)$ |
| 16 | $f(A_{15}, B_{15}, C_{15}, D_{15}, 15)$ | 32 | $g(A_{31}, B_{31}, C_{31}, D_{31}, 12)$ |

Inverting 22 steps

- | | | | |
|----|---|----|---|
| 1 | $f(A_0, B_0, C_0, D_0, 0)$ | 17 | $g(A_{16}, B_{16}, C_{16}, D_{16}, 1)$ |
| 2 | $f(A_1, B_1, C_1, D_1, 1)$ | 18 | $g(A_{17}, B_{17}, C_{17}, D_{17}, 6)$ |
| 3 | $f(A_2, B_2, C_2, D_2, 2)$ | 19 | $g(A_{18}, B_{18}, C_{18}, D_{18}, 11)$ |
| 4 | $f(A_3, B_3, C_3, D_3, 3)$ | 20 | $g(A_{19}, B_{19}, C_{19}, D_{19}, 0)$ |
| 5 | $f(A_4, B_4, C_4, D_4, 4)$ | 21 | $g(A_{20}, B_{20}, C_{20}, D_{20}, 5)$ |
| 6 | $f(A_5, B_5, C_5, D_5, 5)$ | 22 | $g(A_{21}, B_{21}, C_{21}, D_{21}, 10)$ |
| 7 | $f(A_6, B_6, C_6, D_6, 6)$ | | |
| 8 | $f(A_7, B_7, C_7, D_7, 7)$ | | |
| 9 | $f(A_8, B_8, C_8, D_8, 8)$ | | |
| 10 | $f(A_9, B_9, C_9, D_9, 9)$ | | |
| 11 | $f(A_{10}, B_{10}, C_{10}, D_{10}, 10)$ | | |
| 12 | $f(A_{11}, B_{11}, C_{11}, D_{11}, 11)$ | | |
| 13 | $f(A_{12}, B_{12}, C_{12}, D_{12}, 12)$ | | |
| 14 | $f(A_{13}, B_{13}, C_{13}, D_{13}, 13)$ | | |
| 15 | $f(A_{14}, B_{14}, C_{14}, D_{14}, 14)$ | | |
| 16 | $f(A_{15}, B_{15}, C_{15}, D_{15}, 15)$ | | |

Inverting 22 steps

Pick $M_0 \dots M_{11}$

From **1** to **12**, compute $A_{12}B_{12}C_{12}D_{12}$

From **22** to **17**, compute $A_{16}B_{16}C_{16}D_{16}$

Choose M_{12} such that $B_{13} = A_{16}$

Choose M_{13} such that $B_{14} = D_{16}$

Choose M_{14} such that $B_{15} = C_{16}$

Choose M_{15} such that $B_{16} = B_{16}$

Cost: 22 steps

31 steps: same idea...

- | | | | |
|----|---|----|---|
| 1 | $f(A_0, B_0, C_0, D_0, 0)$ | 17 | $g(A_{16}, B_{16}, C_{16}, D_{16}, 1)$ |
| 2 | $f(A_1, B_1, C_1, D_1, 1)$ | 18 | $g(A_{17}, B_{17}, C_{17}, D_{17}, 6)$ |
| 3 | $f(A_2, B_2, C_2, D_2, 2)$ | 19 | $g(A_{18}, B_{18}, C_{18}, D_{18}, 11)$ |
| 4 | $f(A_3, B_3, C_3, D_3, 3)$ | 20 | $g(A_{19}, B_{19}, C_{19}, D_{19}, 0)$ |
| 5 | $f(A_4, B_4, C_4, D_4, 4)$ | 21 | $g(A_{20}, B_{20}, C_{20}, D_{20}, 5)$ |
| 6 | $f(A_5, B_5, C_5, D_5, 5)$ | 22 | $g(A_{21}, B_{21}, C_{21}, D_{21}, 10)$ |
| 7 | $f(A_6, B_6, C_6, D_6, 6)$ | 23 | $g(A_{22}, B_{22}, C_{22}, D_{22}, 15)$ |
| 8 | $f(A_7, B_7, C_7, D_7, 7)$ | 24 | $g(A_{23}, B_{23}, C_{23}, D_{23}, 4)$ |
| 9 | $f(A_8, B_8, C_8, D_8, 8)$ | 25 | $g(A_{24}, B_{24}, C_{24}, D_{24}, 9)$ |
| 10 | $f(A_9, B_9, C_9, D_9, 9)$ | 26 | $g(A_{25}, B_{25}, C_{25}, D_{25}, 14)$ |
| 11 | $f(A_{10}, B_{10}, C_{10}, D_{10}, 10)$ | 27 | $g(A_{26}, B_{26}, C_{26}, D_{26}, 3)$ |
| 12 | $f(A_{11}, B_{11}, C_{11}, D_{11}, 11)$ | 28 | $g(A_{27}, B_{27}, C_{27}, D_{27}, 8)$ |
| 13 | $f(A_{12}, B_{12}, C_{12}, D_{12}, 12)$ | 29 | $g(A_{28}, B_{28}, C_{28}, D_{28}, 13)$ |
| 14 | $f(A_{13}, B_{13}, C_{13}, D_{13}, 13)$ | 30 | $g(A_{29}, B_{29}, C_{29}, D_{29}, 2)$ |
| 15 | $f(A_{14}, B_{14}, C_{14}, D_{14}, 14)$ | 31 | $g(A_{30}, B_{30}, C_{30}, D_{30}, 7)$ |
| 16 | $f(A_{15}, B_{15}, C_{15}, D_{15}, 15)$ | | |

M_{12} input only once...

- | | | | |
|----|---|----|---|
| 1 | $f(A_0, B_0, C_0, D_0, 0)$ | 17 | $g(A_{16}, B_{16}, C_{16}, D_{16}, 1)$ |
| 2 | $f(A_1, B_1, C_1, D_1, 1)$ | 18 | $g(A_{17}, B_{17}, C_{17}, D_{17}, 6)$ |
| 3 | $f(A_2, B_2, C_2, D_2, 2)$ | 19 | $g(A_{18}, B_{18}, C_{18}, D_{18}, 11)$ |
| 4 | $f(A_3, B_3, C_3, D_3, 3)$ | 20 | $g(A_{19}, B_{19}, C_{19}, D_{19}, 0)$ |
| 5 | $f(A_4, B_4, C_4, D_4, 4)$ | 21 | $g(A_{20}, B_{20}, C_{20}, D_{20}, 5)$ |
| 6 | $f(A_5, B_5, C_5, D_5, 5)$ | 22 | $g(A_{21}, B_{21}, C_{21}, D_{21}, 10)$ |
| 7 | $f(A_6, B_6, C_6, D_6, 6)$ | 23 | $g(A_{22}, B_{22}, C_{22}, D_{22}, 15)$ |
| 8 | $f(A_7, B_7, C_7, D_7, 7)$ | 24 | $g(A_{23}, B_{23}, C_{23}, D_{23}, 4)$ |
| 9 | $f(A_8, B_8, C_8, D_8, 8)$ | 25 | $g(A_{24}, B_{24}, C_{24}, D_{24}, 9)$ |
| 10 | $f(A_9, B_9, C_9, D_9, 9)$ | 26 | $g(A_{25}, B_{25}, C_{25}, D_{25}, 14)$ |
| 11 | $f(A_{10}, B_{10}, C_{10}, D_{10}, 10)$ | 27 | $g(A_{26}, B_{26}, C_{26}, D_{26}, 3)$ |
| 12 | $f(A_{11}, B_{11}, C_{11}, D_{11}, 11)$ | 28 | $g(A_{27}, B_{27}, C_{27}, D_{27}, 8)$ |
| 13 | $f(A_{12}, B_{12}, C_{12}, D_{12}, 12)$ | 29 | $g(A_{28}, B_{28}, C_{28}, D_{28}, 13)$ |
| 14 | $f(A_{13}, B_{13}, C_{13}, D_{13}, 13)$ | 30 | $g(A_{29}, B_{29}, C_{29}, D_{29}, 2)$ |
| 15 | $f(A_{14}, B_{14}, C_{14}, D_{14}, 14)$ | 31 | $g(A_{30}, B_{30}, C_{30}, D_{30}, 7)$ |
| 16 | $f(A_{15}, B_{15}, C_{15}, D_{15}, 15)$ | | |

31 steps

Pick $M_0 \dots M_{11}, M_{13}, M_{14}, M_{15}$

From **1** to **12**, compute $A_{12}B_{12}C_{12}D_{12}$

From **22** to **14**, compute $A_{13}B_{13}C_{13}D_{13}$

If $A_{13} = D_{12}$, $C_{13} = B_{12}$, and $D_{13} = C_{12}$:

then choose M_{12} such that $B_{13} = B_{13}$

Cost: $\approx 2^{96} \times 31$ steps

47 steps: M_2 input only twice...

1	$f(\dots, 0)$	17	$g(\dots, 1)$	33	$h(\dots, 5)$
2	$f(\dots, 1)$	18	$g(\dots, 6)$	34	$h(\dots, 8)$
3	$f(\dots, 2)$	19	$g(\dots, 11)$	35	$h(\dots, 11)$
4	$f(\dots, 3)$	20	$g(\dots, 0)$	36	$h(\dots, 14)$
5	$f(\dots, 4)$	21	$g(\dots, 5)$	37	$h(\dots, 1)$
6	$f(\dots, 5)$	22	$g(\dots, 10)$	38	$h(\dots, 4)$
7	$f(\dots, 6)$	23	$g(\dots, 15)$	39	$h(\dots, 7)$
8	$f(\dots, 7)$	24	$g(\dots, 4)$	40	$h(\dots, 10)$
9	$f(\dots, 8)$	25	$g(\dots, 9)$	41	$h(\dots, 13)$
10	$f(\dots, 9)$	26	$g(\dots, 14)$	42	$h(\dots, 0)$
11	$f(\dots, 10)$	27	$g(\dots, 3)$	43	$h(\dots, 3)$
12	$f(\dots, 11)$	28	$g(\dots, 8)$	44	$h(\dots, 6)$
13	$f(\dots, 12)$	29	$g(\dots, 13)$	45	$h(\dots, 9)$
14	$f(\dots, 13)$	30	$g(\dots, 2)$	46	$h(\dots, 12)$
15	$f(\dots, 14)$	31	$g(\dots, 7)$	47	$h(\dots, 15)$
16	$f(\dots, 15)$	32	$g(\dots, 12)$		

Differences propagation, general case

Pick random A_0, B_0, C_0, D_0 and M

- 1 $f(A_0, B_0, C_0, D_0, 0)$
- 2 $f(A_1, B_1, C_1, D_1, 1)$
- 3 $f(A_2, B_2, C_2, D_2, 2)$

Modify C_0 to C_0^*

- X 1 $f(A_0, B_0, C_0^*, D_0, 0)$
- X 2 $f(A_1, B_1, C_1, C_0^*, 1)$
- X 3 $f(C_0^*, B_2, C_2, D_2, 2)$

⇒ all first steps affected (X =state modified)

Difference in C_0 + chosen IV

Pick random A_0, C_0, D_0 and M and set $B_0 = 0$

- 1 $f(A_0, B_0, C_0, D_0, 0)$
- 2 $f(A_1, B_1, C_1, D_1, 1)$
- 3 $f(A_2, B_2, C_2, D_2, 2)$

Modify C_0 to C_0^*

- ✓ 1 $f(A_0, 0, C_0^*, D_0, 0)$
- ✓ 2 $f(A_1, B_1, 0, C_0^*, 1)$
- X 3 $f(C_0^*, B_2, C_2, 0, 2)$

⇒ only step 3 affected

Difference in M_2

Pick random A_0, B_0, C_0, D_0 and M

- 1 $f(A_0, B_0, C_0, D_0, 0)$
- 2 $f(A_1, B_1, C_1, D_1, 1)$
- 3 $f(A_2, B_2, C_2, D_2, 2)$

Modify M_2

- ✓ 1 $f(A_0, B_0, C_0, D_0, 0)$
- ✓ 2 $f(A_1, B_1, C_1, C_0, 1)$
- X 3 $f(A_2, B_2, C_2, D_2, 2)$

⇒ only step 3 affected

Absorbing differences

Pick random A_0, C_0, D_0 and M and set $B_0 = 0$

- 1 $f(A_0, B_0, C_0, D_0, 0)$
- 2 $f(A_1, B_1, C_1, D_1, 1)$
- 3 $f(A_2, B_2, C_2, D_2, 2)$

Modify C_0 to C_0^* and M_2

- ✓ 1 $f(A_0, 0, C_0^*, D_0, 0)$
- ✓ 2 $f(A_1, B_1, 0, C_0^*, 1)$
- ✓ 3 $f(C_0^*, B_2, C_2, 0, 2)$

⇒ nothing changes!

Application to 47-step MD5: key steps

1	$f(\dots, 0)$	17	$g(\dots, 1)$	33	$h(\dots, 5)$
2	$f(\dots, 1)$	18	$g(\dots, 6)$	34	$h(\dots, 8)$
3	$f(\dots, 2)$	19	$g(\dots, 11)$	35	$h(\dots, 11)$
4	$f(\dots, 3)$	20	$g(\dots, 0)$	36	$h(\dots, 14)$
5	$f(\dots, 4)$	21	$g(\dots, 5)$	37	$h(\dots, 1)$
6	$f(\dots, 5)$	22	$g(\dots, 10)$	38	$h(\dots, 4)$
7	$f(\dots, 6)$	23	$g(\dots, 15)$	39	$h(\dots, 7)$
8	$f(\dots, 7)$	24	$g(\dots, 4)$	40	$h(\dots, 10)$
9	$f(\dots, 8)$	25	$g(\dots, 9)$	41	$h(\dots, 13)$
10	$f(\dots, 9)$	26	$g(\dots, 14)$	42	$h(\dots, 0)$
11	$f(\dots, 10)$	27	$g(\dots, 3)$	43	$h(\dots, 3)$
12	$f(\dots, 11)$	28	$g(\dots, 8)$	44	$h(\dots, 6)$
13	$f(\dots, 12)$	29	$g(\dots, 13)$	45	$h(\dots, 9)$
14	$f(\dots, 13)$	30	$g(\dots, 2)$	46	$h(\dots, 12)$
15	$f(\dots, 14)$	31	$g(\dots, 7)$	47	$h(\dots, 15)$
16	$f(\dots, 15)$	32	$g(\dots, 12)$		

The attack

Stage 1: MITM

Pick M and IV with $B_0 = 0$,

1. store $(A_{29}, B_{29}, C_{29}, D_{29})$ for all $2^{32} C_0$'s (forward)
2. store $(A_{30}, B_{30}, C_{30}, D_{30})$ for all $2^{32} C_{47}$'s (backward)

Find entries such that

$$A_{30} = D_{29}$$

$$D_{30} = C_{29}$$

$$C_{30} = B_{29}$$

\equiv 96-bit equality; 2^{64} choices \Rightarrow repeat 2^{32} times

The attack

Stage 2: correction

Modify M_2 such that

$$B_{30} = g(A_{29}, B_{29}, C_{29}, D_{29}, 2)$$

Modify C_0 accordingly

⇒ 96-bit preimage ($C_0 + C_{47}$ is random) with prob. 2^{-32}

Total cost: 2^{96} trials for a 128-bit preimage

Summary

Preimages for MD5's **compression function** with

- ▶ chosen message except M_1 and M_2
- ▶ IV with $B_0 = 0$ and random C_0
- ▶ storage for 2^{36} bytes (64 Gb)
- ▶ 2^{96} compressions

By bruteforce:

- ▶ random message
- ▶ chosen IV
- ▶ negligible memory
- ▶ 2^{128} compressions

HAVAL

- ▶ 1992: publication (Zheng, Pieprzyk, Seberry)
- ▶ 2003: **collision** attack (3-pass)
- ▶ 2006: **collision** attack (4- and 5-pass)
- ▶ 2008: (partial) **second-preimage** attack (3-pass)

Function similar to MD5 with

- ▶ 256-bit chain values
- ▶ 1024-bit blocks
- ▶ 3, 4, or 5 rounds

Preimages for 3-pass HAVAL

Same strategy as for MD5

- ▶ identify **absorption** properties in the initial steps
 - ▶ MITM
 - ▶ modify a M_i to complete the MITM
 - ▶ correct initial steps
- ⇒ 2 attacks in 2^{224} and storage of 2^{69} bytes
(vs. 2^{256} and negligible memory)

Extension to the iterated hash

Attacks presented for the compression function
(with IV partially random, no padding)

Restrictions for the hash function:

- ▶ **padding**: not a problem, because message chosen
- ▶ **fixed IV**: makes direct application impossible

Iterated hash: basic MITM

Given image \mathbf{H} :

1. compute a **list of images** from the fixed IV

$$(M_i, \text{compress}(\text{IV}, M_i))_i$$

2. compute a **list of preimages**

$$(H_i, M'_i)_i, \text{ compress}(H_i, M'_i) = \mathbf{H}$$

Find entries such that

$$\text{compress}(\text{IV}, M_i) = H_j$$

Cost: 2^{113} trials + 2^{36} bytes for MD5, $2^{241} + 2^{69}$ for HAVAL

Iterated hash: tree technique

Mendel & Rijmen (ICISC'07), Leurent (FSE'08)

Build a tree using **multi-target preimages**

Cost: 2^{102} trials + 2^{39} bytes for MD5, $2^{230} + 2^{71}$ for HAVAL

Conclusion

	compressions	bytes
47-step MD5	2^{102}	2^{39}
3-pass HAVAL	2^{230}	2^{71}

Independent work by Sasaki & Aoki @ACISP 2008:
44-step MD5 (starting from step 6) in 2^{96}

Questions?

Is it effectively faster than bruteforce?

→ Arguably yes (but not 2^{26} times faster)

Same strategy applies to MD4?

→ No (because no M_i at very start and very end)

Same strategy applies to SHA-0/1/2?

→ No (nontrivial message expansion), cf. next talk